

E-ID Hands-on Workshop

Keeping identities safe and sound





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Agenda

- 9:15am welcome coffee
- 9:30am Overview of E-ID Landscape Dr. Imad Aad
- 11:30am lunch
- 1:00pm Hands-on training, led by C4DT:
- Starting with signed verified credentials and selective disclosure
- Using BBS+ to add unlinkability
- Introducing zero-knowledge proofs to reduce information leakage
- 4:30pm Wrap-up: Lessons learned and how to go forward



C4DT - TLDR







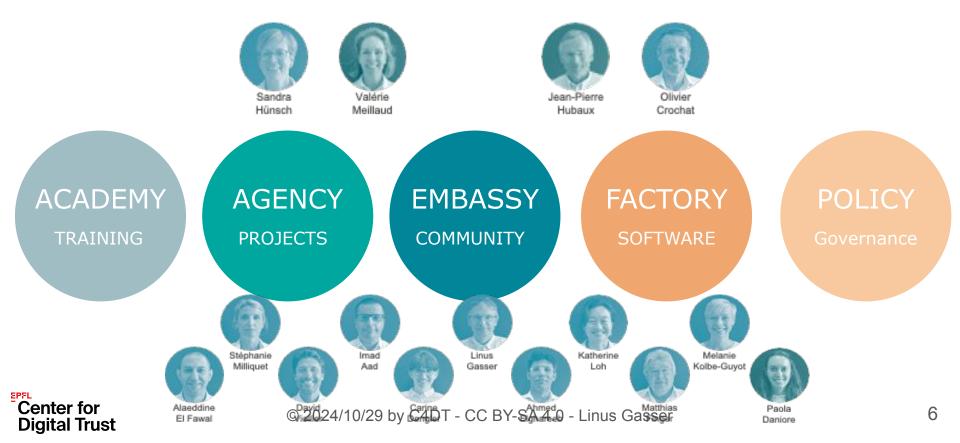
5 Domains



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Team of 15 (+2) People





C4DT Factory - Overview



INCUBATOR

- Project Presentation
- Paper -> Real world
- Demonstrators / Market



KNOWLEDGE

- Explore subjects
- Article / Blog posts
- Conferences / Teaching

HANDS-ON WORKSHOPS

- 1-day trainings
- Share latest research
- Real-world input



COMMUNITY

- Research Software Engineers
- EPFL Labs Factory -Partners
- Conferences

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C4DT Factory in October / November 2024







Curtain Call for our Demonstrators

C4DT Factory Update 2024/10 Nov. 1st Deepfake round-table and workshop Nov. 19th & 26th

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Overview of E-ID Landscape

Dr. Imad Aad, C4DT





E-ID Hands-on Workshop

Keeping identities safe and sound





Program

- 1. Signing simply with RSA
- 2. Unlinkable proofs using BBS+
- 3. Predicate proofs with ZKPs
- 4. ZKP Considerations

For subjects 1-3:

- 1. Short theory
- 2. Jupyter exercises
- 3. Discussion
- 4. Longer coding exercise

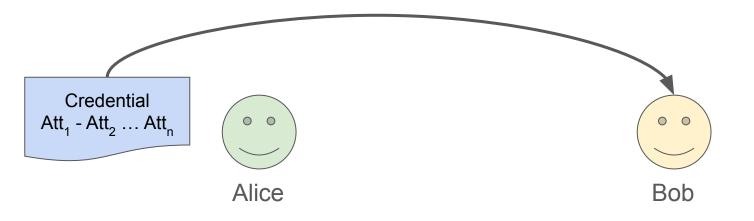


1 - Signing Simply with RSA





Attribute Sharing

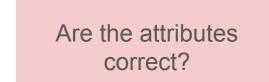


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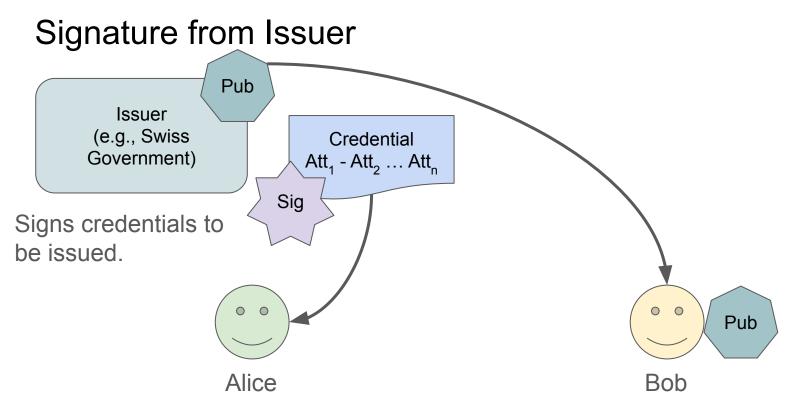
Attribute Sharing - 1st Problem





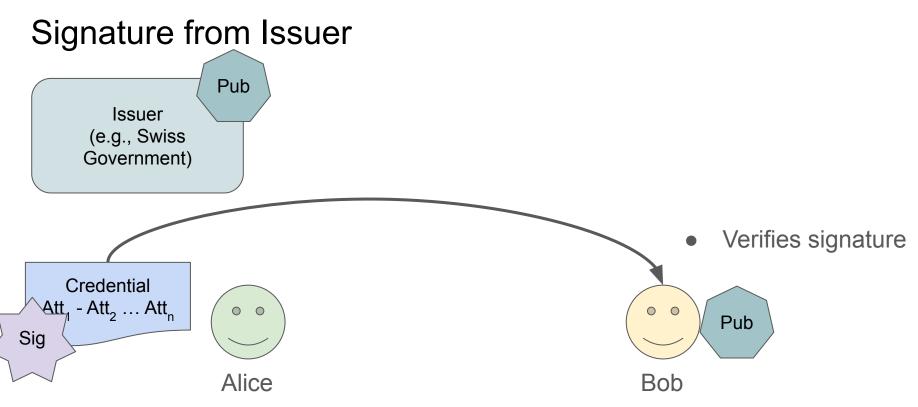








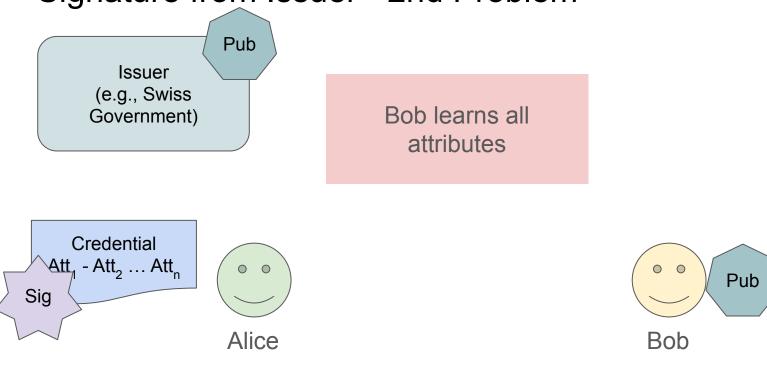




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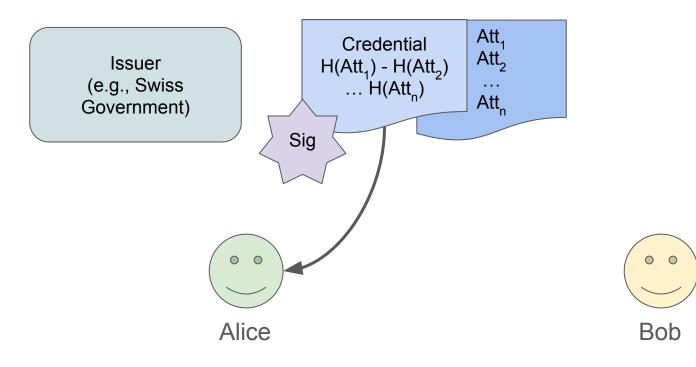
Signature from Issuer - 2nd Problem







Selective Disclosure

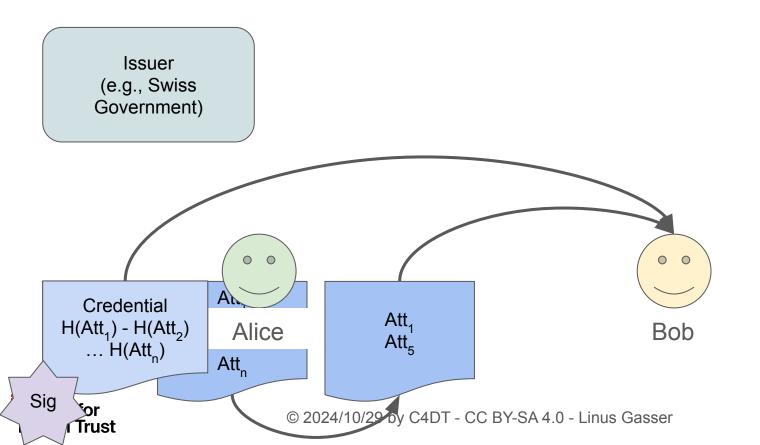




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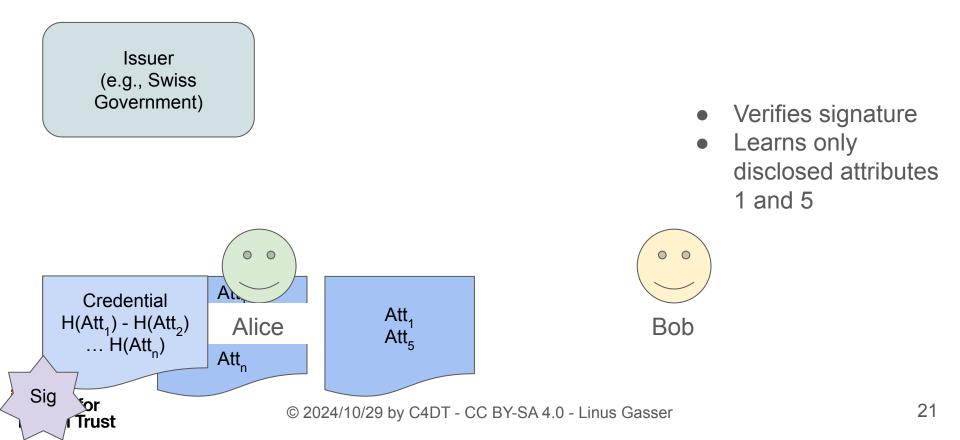


Selective Disclosure



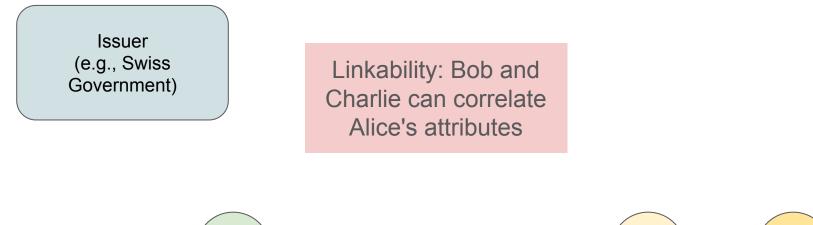


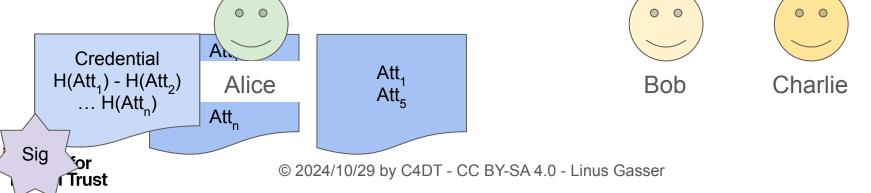
Selective Disclosure





Selective Disclosure - 3rd Problem





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Exercise 1 - Signing Simply with RSA





Wrap-up slide

- The issuer allows the verifier to trust the data from the holder
- Selective disclosure can hide personal data to the verifier
- For low-entropy data, even cryptographic hashes do not provide anonymity
- LD-JSON Verified Credentials from EU Digital Wallet are linkable



2 - Unlinkable proofs using BBS+



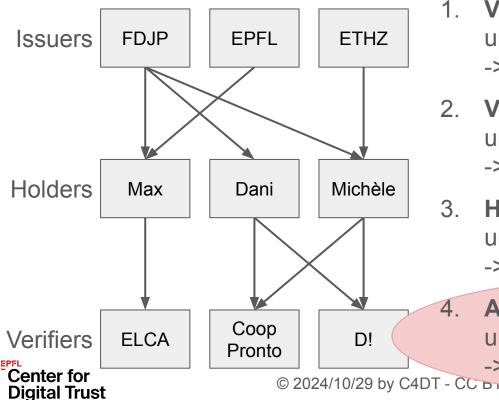


Why Unlinkability?

- No correlation between visits
- Reduces attack surface if data leaks
- Privacy / Profiling
 - less knowledge about visitors -> less influence
 - no following of holders -> physical security (e.g., stalkers)



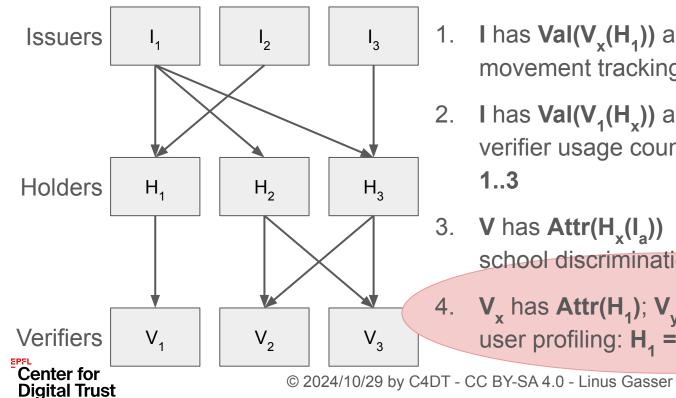
Unlinkability Vows (in addition to anonymity)



- Validity check by Coop and D! on Dani unlinkable by the FDJP
 -> movement tracking
- 2. Validity check by D! on Dani and Michèle unlinkable by the FDJP
 -> counting of usage by a verifier
 - Has CH Master Degree check by ELCA unlinkable to EPFL or ETHZ
 -> discrimination against a school
 - Age check by Coop and D! on Dani unlinkable by Coop and D! -> user profiling



Unlinkability Vows (in addition to anonymity)



I has $Val(V_x(H_1))$ and $Val(V_v(H_2))$ movement tracking: $H_1 = ? H_2 \forall x, y \in 1..3$

2. I has $Val(V_1(H_y))$ and $Val(V_2(H_y))$ verifier usage counting: $V_1 = ?V_2 \forall x, y \in$

school discrimination: $a = ?2,3 \forall x \in 1..3$

4. V_x has Attr(H_1); V_y has Attr(H_2) user profiling: $H_1 = ? H_2 \forall x, y \in 1..3$

How to Make it Unlinkable

- 1. and 2. validity or revocation check
- Cryptographic accumulators slow and potentially huge
- 3. Issuer hiding
 - Create "meta issuer" issuer of issuers
- 4. User profiling
 - BBS+ signatures



Avoid User Profiling with BBS+

If V_x has Attr(H_1); V_y has Attr(H_2), it's difficult to verify if $H_1 = H_2$, $\forall x, y \in 1..3$

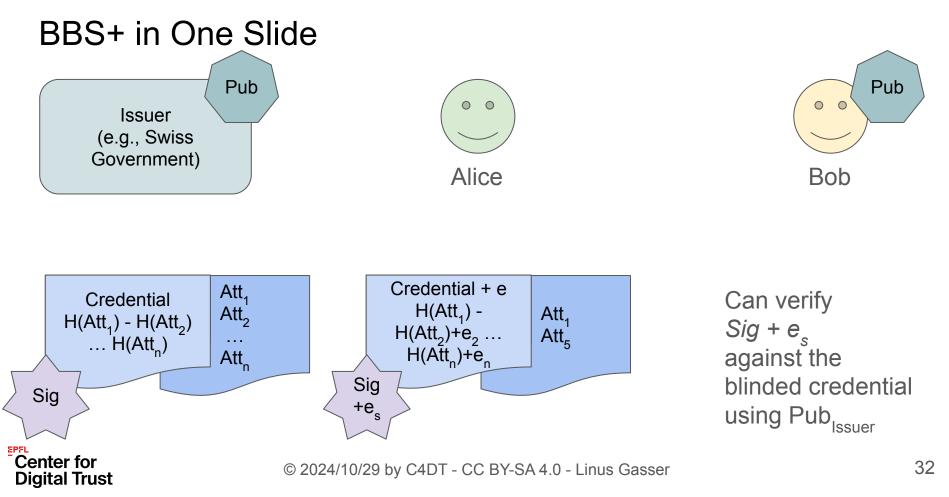
- Issuer signature needs to be blinded (valid but different each time)
- Hashes of the non-disclosed fields need to be blinded
- BBS(+) to the rescue
 - Zero-knowledge proof:

Here is a proof that I know a signature of the following hash(es)

- BBS: original paper, security proof only later
- BBS+: added a random factor to create a security proof
- BBS#: extension proposed by Orange to do holder binding
- Short BBS: not using pairing-based cryptography

Blinding disclosed fields -> Predicate Zero Knowledge Proofs, not in BBS+!







Exercise 2 - Unlinkable proofs using BBS+



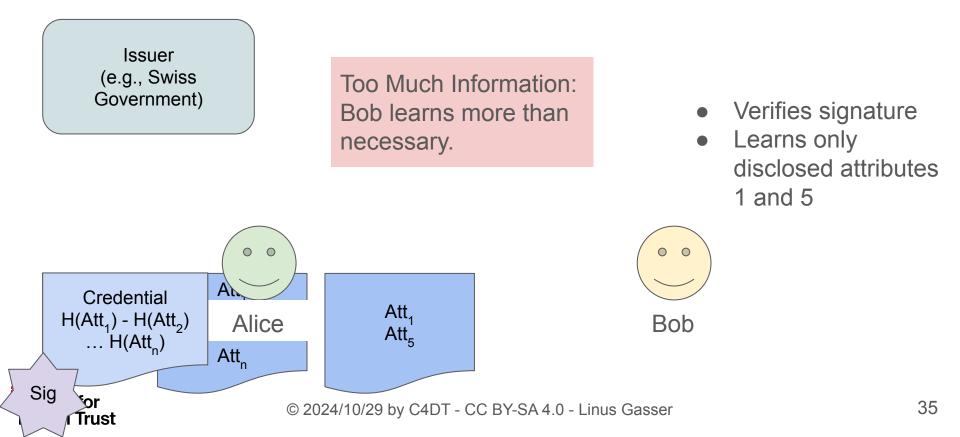


What we Learnt

- BBS+ creates unlinkable proofs
- It can selectively disclose fields chosen by the holder
- Hover, the disclosed fields might still be used to link proofs



Selective Disclosure - 4th Problem





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Unlinkability - and Now?

Disclosed values are fully visible, for example

- Birthdate (when you only want to prove you're > 65)
- Salary (instead of proving you earn less than 30k)
- Address (reduction for a ticket bc you live in VD)

This is not desirable because of:

- Privacy: you don't want to give away that data
- De-anonymization: when combining fields, you can get a very small anonymity set (male, 1.1.1978, 1015)



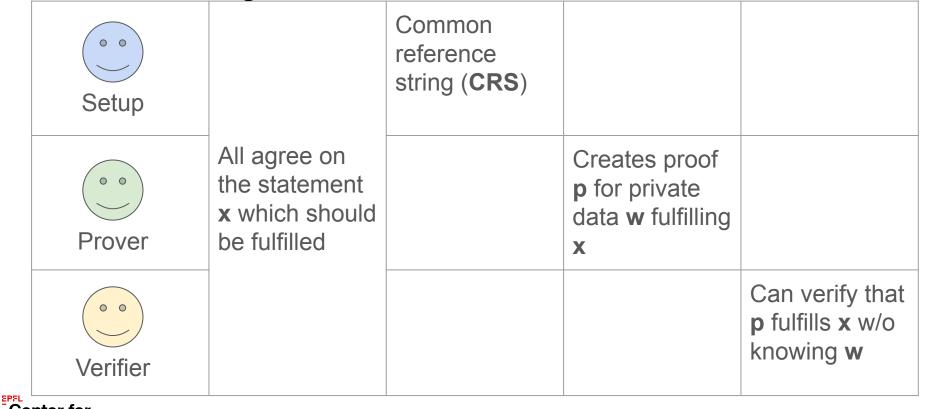


3 - Predicate Proofs with ZKPs





Zero Knowledge Proofs 101



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An Example of a Statement

Wanting to buy a ticket with a reduction for retired people:

Proving the issuer signed a verified credential which includes an age \geq = 65:

- All agree on the condition **x**:
 - I know a signature Sig_{issuer}+e_{sig} to a hash H_A+e_A verifiable by Pub_{issuer} AND
 - I know a number N_A which hashes to $H_A + e_A$ AND
 - \circ **N**_A is above or equal to 65
- The holder creates a proof **p** for **x** using their **w**
- The verifier can check **p** fulfills **x**, knowing only **Pub**_{issuer}



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Biggest Zero Knowledge Proof Families in 2024

Name	Foundation	Setup	Proof creation	Verification
SNARK	Bilinear pairings, elliptic curves PQ: No	Yes Time: long	Size: constant Time: fast (w/o setup)	Time: fast
STARK	Hash functions PQ: Yes	No	Size: large Time: slow	Time: fast
Bulletproofs	Elliptic curves PQ: No	No	Size: medium Time: slow	Time: medium

2024/10 - depends also on complexity of statement **x**

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Some Zero Knowledge Terms

- **Completeness**: If the statement is true, an honest prover will be able to convince an honest verifier of this fact.
- **Soundness**: If the statement is false, no dishonest prover can convince an honest verifier that it is true, except with a very small probability.
- **Zero-Knowledge**: If the statement is true, the verifier learns nothing other than the fact that the statement is true.
- **Interactive**: the verifier interacts over many rounds with the prover, until they are convinced of the statement. Sigma protocols are interactive ZKPs.
- **Succinctness**: the proof size should be small, and the verification time should be fast



Exercise 3 - Predicate proofs with ZKPs





Wrap-up slide

The good:

- Zero Knowledge Proofs allow to minimize the data leakage from the credentials
- The docknetwork/crypto library has a very powerful mechanism to set up a ZKP statement

The bad:

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- There are no standards yet it is very new
- Some statements are still very complicated to express



4 - ZKP Considerations



Difference Between ZKP Systems

- Setup: either with (zkSNARK) or without (zkSTARK, Bulletproofs)
 - with: smaller and faster proofs and verifications, but need to trust the setup
 - without: no trust needed
 - o as seen in the exercises, fast advancing research turns the tables
- Statement complexity
- Setup: time and size ms to seconds; 1-100kB
- Proof creation: time and size ms to minutes; 100B to xMB
- Verification: time ms to seconds



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(Lego)Groth16 <-> Bulletproofs++

- Groth16 is an "old" algorithm which is well understood
- Bulletproofs(++) is more advanced, and looks like it could replace Lego16
- LegoGroth16 is an example of combining various ZKP algorithms
- The docknetwork/crypto library adds yet another layer

Comparison in exercise:

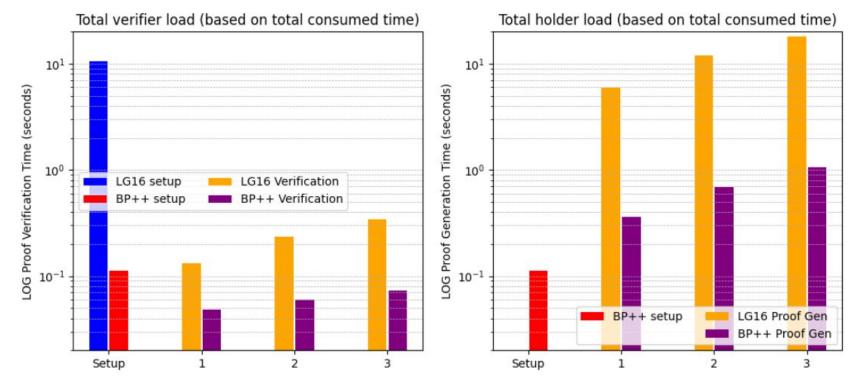
- Computation cost:
 - Server: setup and verify
 - Client: setup and create proof
- Communication cost:
 - Server -> client: setup material
 - Client -> server: proof



Exercise 4 - ZKP Considerations

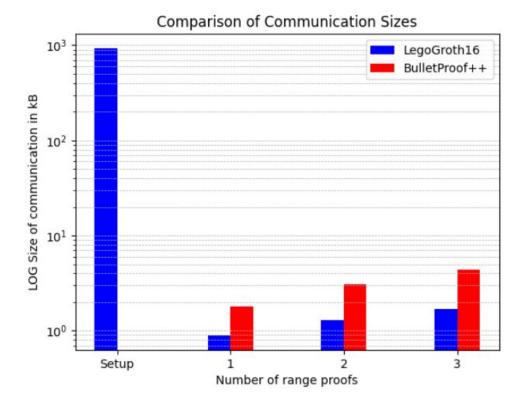


Setup and Proof Generation - Logarithmic y-scale!



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Communication Sizes



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Interpretation

This is very specific to the *docknetwork/crypto* library:

- Special setup to create composed proofs
- Not optimized for 'simple' range proofs

Generally:

- The setup for the LegoGroth16 can be re-used by the verifier
- The setup for Bulletproofs++ must be done every time
- The communication size for LegoGroth16 is very high





Conclusions



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Setting up a Trustworthy E-ID

- What is important?
 - Convince Swiss citizens that E-ID is trustworthy
 - Use Cases for the E-ID
- Questions for the Swiss E-ID
 - ZKP for ECDSA signatures for holder binding
 - Which basic signatures scheme to use
- Standardizations
 - BBS+ has an IETF draft
 - Nothing yet for ZKPs